Problem 1 Let a be an odd positive integer. Prove that if $d \mid (a^2 + 2)$ then $d \equiv 1 \pmod{8}$ or $d \equiv 3 \pmod{8}$. **Solution** If $p \mid (a^2 + 2)$ and p is the prime, then -2 is a quadratic residue modulo p. It follows that for the Legendre symbol $\left(\frac{-2}{p}\right) = 1$, where $\left(\frac{1}{2}\right)$ is the Legendre symbol. Using the properties of Legendre symbol we obtain

$$\left(\frac{-2}{p}\right) = \left(\frac{-1}{p}\right) \cdot \left(\frac{2}{p}\right) = (-1)^{\frac{p-1}{2} + \frac{p^2-1}{8}}$$

Thus the number $\frac{p-1}{2} + \frac{p^2-1}{8}$ is even and so

$$p^2 + 4p \equiv 5 \pmod{16}.\tag{1}$$

On the other hand p is an odd prime, thus p is of the form 8k + 1, 8k + 3, 8k + 5 or 8k + 7. This and (1) yield that the prime p is of the form 8k + 1 or 8k + 3. The product of an even number of numbers of the form 8k + 3 is a number b fulfilling $b \equiv 1 \pmod{8}$, and the product of an odd number of numbers of the form 8k + 3 is c fulfilling $c \equiv 3 \pmod{8}$. The product e of numbers of the form 8k + 1 satisfies $e \equiv 1 \pmod{8}$. This yields the assertion.

Problem 2 Let $\alpha \in (0,1]$ be a given real number and let the real sequence $\{a_n\}_{n=1}^{\infty}$ satisfy the inequality

$$a_{n+1} \le \alpha a_n + (1-\alpha)a_{n-1}$$
 for $n = 2, 3, \dots$

If $\{a_n\}$ is bounded prove that then it must be convergent.

Solution Since (a_n) is bounded there exists both $\liminf a_n = l$ and $\limsup a_n = L$. We shall prove that L = l. Let suppose the contrary L > l. By the definition of lim sup and lim inf we know that there exist subsequencies (a_{n_k}) and (a_{m_k}) of (a_n) which converge to l and L.

Since $L = \limsup a_n$ we have:

$$(\forall \varepsilon > 0) (\exists n_0 \in \mathbb{N}) (\forall n \ge n_0) a_n < L + \varepsilon.$$

On the other hand

$$(\forall \varepsilon > 0) (\exists n_1 \in \mathbb{N}) (\forall n \ge n_1) l - \varepsilon < a_n$$

We choose m_{k_0} such that: $m_{k_0} > n_1, l - \varepsilon < a_{m_{k_0}} < l + \varepsilon$. By the inequality we obtain:

$$a_{m_{k_0+1}} \le (1-\alpha)a_{m_{k_0}} + \alpha a_{m_{k_0-1}} < (1-\alpha)(l+\varepsilon) + \alpha(L+\varepsilon) = \varepsilon + (1-\alpha)l + \alpha L$$

and

$$a_{m_{k_0+2}} \le (1-\alpha)a_{m_{k_0+1}} + \alpha a_{m_{k_0}} < (1-\alpha)(\varepsilon + (1-\alpha)l + \alpha L) + \alpha(l+\varepsilon) = \varepsilon + ((1-\alpha)^2 + \alpha)l + (1-\alpha)\alpha L$$

It is not difficult to demonstrate that we can choose ε such that:

$$\begin{split} \varepsilon + (1-\alpha) l + \alpha L &< L - \frac{\varepsilon}{2} \,, \\ \varepsilon + (1-\alpha+\alpha^2) l + (\alpha-\alpha^2) L &< L - \frac{\varepsilon}{2} \end{split}$$

using the functions: $f(x) = \varepsilon + (1-x)l + xL$ and $g(x) = \varepsilon + (1-x-x^2)l + (x-x^2)L$, $x \in [0,1]$. From that we have: $a_{m_{k_0+1}} \leq L - \frac{\varepsilon}{2}$, $a_{m_{k_0+2}} \leq L - \frac{\varepsilon}{2}$. From that we obtain:

$$a_{m_{k_0+3}} \le (1-\alpha)a_{m_{k_0+1}} + a_{m_{k_0}} < (1-\alpha)\left(L - \frac{\varepsilon}{2}\right) + \alpha\left(L - \frac{\varepsilon}{2}\right) = L - \frac{\varepsilon}{2}$$

therefore by induction we get:

$$a_n < L - \frac{\varepsilon}{2}, \quad \forall n \ge m_{k_0+1}.$$

The last formula yields:

$$\limsup a_n \le L - \frac{\varepsilon}{2} < L,$$

which is a contradiction. Thus L = l.

Problem 3 Let c_1, c_2, \ldots, c_n be real numbers such that

$$c_1^k + c_2^k + \dots + c_n^k > 0 \quad \text{for } k = 1, 2, \dots$$
 (1)

Let us put

$$f(x) = \frac{1}{(1 - c_1 x)(1 - c_2 x)\dots(1 - c_n x)}.$$

Show that $f^{(k)}(0) > 0$ for all k = 1, 2, ...Solution Put $g(x) := \log f(x)$. Then

$$g(x) = -\sum_{j=1}^{n} \log(1 - c_j x)$$
 and $g^{(k)}(x) = \sum_{j=1}^{n} \frac{(k-1)! c_j^k}{(1 - c_j x)^k}.$

Assumption (1) and the above result give

$$g^{(k)}(0) = c_1^k + c_2^k + \dots + c_n^k > 0.$$

Now observe that if all the derivatives of g at the origin are positive, then e^g has the same property. For this show by induction that

$$(\mathbf{e}^g)^{(k)} = \mathbf{e}^g \cdot S$$

where S is a finite sum of terms of the form

$$a(g^{(l_1)})^{m_1}(g^{(l_2)})^{m_2}\dots(g^{(l_r)})^{m_r},$$

where a, l_j and m_j are positive integers. For instance:

$$\begin{aligned} (e^{g(x)})' &= e^{g(x)}g'(x) \\ (e^{g(x)})'' &= e^{g(x)}((g'(x))^2 + g''(x)) \\ (e^{g(x)})''' &= e^{g(x)}((g'(x))^3 + 3g'(x)g''(x) + g'''(x)) \\ (e^{g(x)})^{(4)} &= e^{g(x)}((g'(x))^4 + 6(g'(x))^2g''(x) + 3(g''(x))^2 + 4g'(x)g'''(x) + g^{(4)}(x)) \end{aligned}$$

Problem 4-M Find all real numbers a > 0 for which the series

$$\sum_{n=1}^{\infty} \frac{a^{f(n)}}{n^2}$$

is convergent, where f(n) denotes the number of zeros in the decimal expansion of n. Solution For n = 0, 1, ... let us consider k's fulfilling the inequality:

$$10^n \le k < 10^{n+1}.\tag{1}$$

For a fixed $0 \leq j \leq n$ there are

$$\binom{n}{j}9^{n-j+1}$$
 k's fulfilling (1) such that $f(k) = j$.

Consequently,

$$\frac{9(9+a)^n}{10^{2n+2}} = \frac{9}{10^{2n+2}} \left(\sum_{j=0}^n \binom{n}{j} a^j 9^{n-j} \right) \le \sum_{k=10^n}^{10^{n+1}-1} \frac{a^{f(k)}}{k^2} \le \frac{9}{10^{2n}} \left(\sum_{j=0}^n \binom{n}{j} a^j 9^{n-j} \right) = \frac{9(9+a)^n}{10^{2n}} \,.$$

This implies that

$$\frac{9}{100}\sum_{n=0}^{\infty} \left(\frac{9+a}{100}\right)^n \le \sum_{n=1}^{\infty} \frac{a^{f(n)}}{n^2} \le 9\sum_{n=0}^{\infty} \left(\frac{9+a}{100}\right)^n.$$

From the last inequalities we conclude that our series is convergent exactly when

$$\frac{9+a}{100} < 1$$

0 < a < 91.

or

	l
_	l

Problem 4-I Let us have declared:

```
const N_MAX = 255;
type
   tR = array[1..N_MAX] of real;
   tN = array[1..N_MAX] of integer;
```

and function random without parameters which returns real random values distributed uniformly in [0, 1). You need to choose K integer numbers $(1 < K < N, N MAX \ge N)$ without repetitions under the condition

that the probability of the choice of a number *i* equals a given P_i , $\sum_{i=1}^{N} P_i = 1$. Write the procedure in Pascal that returns such K integer numbers in the first K elements of the vector of tN type. Input parameters of the procedure are K, N and the vector of P_i .

Solution

Problem 1 Decide whether it is possible to cover a 3-dimensional Euclidean space with lines which are pairwise skew (i.e. not coplanar).

Solution There is (for example) the following covering:

lines
$$p_{a,b} = \{A_{a,b} = [a, b, 0]; s_{a,b} = (-b, a, 1)\},\$$

where $A_{a,b}$ is point and $s_{a,b}$ is vector of $p_{a,b}$. We show that

- 1. if $[a, b] \neq [c, d]$ then $p_{a,b} \cap p_{c,d} = \emptyset$; and
- 2. for each X = [x, y, z] there is [a, b] such that $X \in p_{a,b}$.
- 1. Let $[a, b] \neq [c, d]$ and $p_{a,b} \cap p_{c,d} = \emptyset$. Then (from the parametric expression)

$$a - tb = c - rd$$
$$b + ta = d + rc$$
$$t = r$$

for some real t and r. From the third equality t = r and from the first and the second we have a-c = t(b-d)and b-d = -t(a-c). By linear combination (first times (a-c) plus second times (b-d)) we have

$$(a-c)^{2} + (b-d)^{2} = 0.$$

But this contradicts $[a, b] \neq [c, d]$.

2. For X = [x, y, z] put $a = \frac{x+yz}{1+z^2}$; $b = \frac{y-xz}{1+z^2}$. Then $X = A_{a,b} + z \cdot s_{a,b}$.

This implies that it is possible to cover 3-dimensional Euclidean space with lines which are pairwise skew. \Box

Problem 2 Let $f: \mathbb{C} \to \mathbb{C}$ be a holomorphic function such that |f(z)| = 1 for all $z \in \mathbb{C}$ with |z| = 1. Prove that there are $\theta \in \mathbb{R}, k \in \{0, 1, 2, ...\}$ such that

$$f(z) = e^{i\theta} z^k$$
 for any $z \in \mathbb{C}$.

Solution We know that

$$|f(z)| = 1$$
 for all $|z| = 1$. (1)

Let $\Delta := \{z \in \mathbb{C} : |z| < 1\}$. Let $\{z_1, \ldots, z_n\}$ be zeros (counted with multiplicity) of the function f belonging to Δ . The number of zeros is finite because of the identity principle and (1).

Let us define the following function:

$$g(z) := \frac{f(z)}{\prod_{j=1}^{n} \frac{z-z_j}{1-\bar{z}_j z}}, \quad \text{for all } z \text{ with } z\bar{z}_j \neq 1.$$

From the choice of z_j we certainly have that g is holomorphic in some neighbourhood of $\overline{\Delta}$ (even in $\mathbb{C} \setminus \{1/\overline{z}_1, \ldots, 1/\overline{z}_n\}$).

Because of (1) and the fact that

$$\left|\frac{z-z_j}{1-\bar{z}_j z}\right| = 1$$
 for all $|z| = 1$

we have that

|g(z)| = 1 for all |z| = 1. (2)

Moreover,

$$g(z) \neq 0 \quad \text{for all } z \in \Delta.$$
 (3)

In view of (2) and (3), the minimum and maximum principles applied to the mapping g imply that

$$|g| \equiv 1 \text{ on } \Delta,$$

which gives us that

$$g \equiv e^{i\theta}$$
 for some $\theta \in \mathbb{R}$.

Therefore,

$$f(z) = e^{i\theta} \prod_{j=1}^n \frac{z-z_j}{1-\bar{z}_j z}, \quad z \in \mathbb{C}.$$
 And the last inequality is possible iff $z_1 = \cdots = z_n = 0.$

Problem 3 Let $u \in C^2(\overline{D}), u = 0$ on ∂D , where D is an open unit ball in \mathbb{R}^3 . Prove that the following inequality

$$\int_{D} |\operatorname{grad} u|^2 \, \mathrm{d}V \le \varepsilon \int_{D} (\Delta u)^2 \, \mathrm{d}V + \frac{1}{4\varepsilon} \int_{D} u^2 \, \mathrm{d}V$$

holds for all $\varepsilon > 0$.

Solution Since the inequality must hold for every $\varepsilon > 0$ it will hold too for those ε for which the function $g(\varepsilon) = \varepsilon \int_D (\Delta u)^2 dV + \frac{1}{4\varepsilon} \int_D u^2 dV$ takes a minimum, i.e., for $\varepsilon = \frac{1}{2} \sqrt{\frac{I_u}{I_l}} (g'(\varepsilon) = I_l - \frac{1}{4\varepsilon^2} I_u = 0)$. That minimum is $g(\frac{1}{2}\sqrt{\frac{I_u}{I_l}}) = \sqrt{I_u I_l}$. Therefore it suffices to show that

$$\int_{D} |\operatorname{grad} u|^2 \, \mathrm{d} V \le \sqrt{\int_{D} (\Delta u)^2 \, \mathrm{d} V \int_{D} u^2 \, \mathrm{d} V} \,.$$

By placing $(P, Q, R) = u(u_x, u_y, u_z)$ in the Gauss-Ostrogradski formula we get:

$$\int_{D} |\operatorname{grad} u|^2 \, \mathrm{d}V + \int_{D} u \Delta u \, \mathrm{d}V = -\int_{\partial D} u \frac{\partial u}{\partial n} \, \mathrm{d}S = 0,$$

where n is inner normal. From this by the Cauchy-Schwartz inequality and the conditions of that statement we obtain

$$\int_{D} |\operatorname{grad} u|^2 \, \mathrm{d}V = -\int_{D} u \Delta u \, \mathrm{d}V \le \sqrt{\int_{D} (\Delta u)^2 \, \mathrm{d}V \int_{D} u^2 \, \mathrm{d}V} \,.$$

Problem 4-M Prove that

$$\sum_{n=1}^{\infty} \frac{n^2}{(7n)!} = \frac{1}{7^3} \sum_{k=1}^{2} \sum_{j=0}^{6} e^{\cos\left(\frac{2\pi j}{7}\right)} \cdot \cos\left(\sin\left(\frac{2\pi j}{7}\right) + \left(\frac{2\pi jk}{7}\right)\right).$$

Solution

$$\sum_{n=1}^{\infty} \frac{n^2}{(7n)!} = \frac{1}{49} \sum_{n=1}^{\infty} \frac{7n(7n-1)+7n}{(7n)!} = \sum_{k=1}^{2} \sum_{n=1}^{\infty} \frac{1}{(7n-k)!}$$

We have

$$A = \sum_{k=1}^{2} \sum_{j=0}^{6} e^{\frac{2\pi i j}{7} + \frac{2\pi i j k}{7}} = \sum_{k=1}^{2} \sum_{j=0}^{6} e^{\frac{2\pi i j k}{7}} e^{e^{\frac{2\pi i j k}{7}}} = \sum_{k=1}^{2} \sum_{j=0}^{6} e^{\frac{2\pi i j k}{7}} \sum_{n=0}^{\infty} \frac{e^{\frac{2\pi i j k}{7}}}{n!} = \sum_{k=1}^{2} \sum_{n=0}^{\infty} \frac{1}{n!} \sum_{j=0}^{6} e^{\frac{2\pi i j (n+k)}{7}}$$

If $7 \mid N$, then

$$\sum_{j=0}^{6} e^{\frac{2\pi i j N}{7}} = 7.$$

If $7 \nmid N$, then

$$\sum_{j=0}^{6} e^{\frac{2\pi i j N}{7}} = \frac{e^{2\pi i N} - 1}{e^{\frac{2\pi i j N}{7}} - 1} = 0.$$

It follows that

$$A = 7 \sum_{k=1}^{2} \sum_{n=1}^{\infty} \frac{1}{(7n-k)!} \,.$$

Thus

$$\begin{split} \sum_{n=1}^{\infty} \frac{n^2}{(7n)!} &= \frac{1}{7^3} \sum_{k=1}^2 \sum_{j=0}^6 e^{\frac{2\pi i j}{7} + \frac{2\pi i j k}{7}} = \frac{1}{7^3} \sum_{k=1}^2 \sum_{j=0}^6 e^{\cos(\frac{2\pi j}{7}) + i(\sin(\frac{2\pi j}{7}) + (\frac{2\pi j k}{7}))} = \\ &= \frac{1}{7^3} \sum_{k=1}^2 \sum_{j=0}^6 e^{\cos(\frac{2\pi j}{7})} \cdot \cos\left(\sin\left(\frac{2\pi j}{7}\right) + \left(\frac{2\pi j k}{7}\right)\right). \end{split}$$

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Problem 4-I Problem Div_3 is specified as follows:

Input: any program P,

Output: a finite set D(P) of strings of 0's and 1's,

where it holds that the program P solves problem Div_3 iff it outputs the correct answers for inputs from D(P).

Theorem For any program TRANSF which transforms programs in some way (i.e. for any given program P it constructs some program P', denoted by P' = TRANSF(P)) there is a program P_0 whose input/output behaviour is not changed by the transformation (i.e. P_0 and $TRANSF(P_0)$ yield the same outputs for the same inputs).

Solution Suppose there is some GEN-TEST-DATA with the property described above. Now construct a program TRANSF which works as follows:

When given a program P, it constructs D(P) by using GEN-TEST-DATA and then constructs P' whose behaviour can be described as follows:

```
\begin{array}{l} \texttt{s:= input;} \\ \texttt{if } s \in D(P)\texttt{;} \\ \texttt{then output YES or NO according to whether or not s is} \\ \texttt{the binary code of a number divisible by 3;} \\ \texttt{else output YES;} \end{array}
```

Due to the Recursion Theorem there is a program P_0 s.t. its input/output behaviour is the same as the behaviour of $TRANSF(P_0)$, and it can be described as follows:

```
\label{eq:s:input} \begin{array}{l} \texttt{s:= input;} \\ \texttt{if } s \in D(P_0)\texttt{;} \\ \texttt{then output YES or NO according to whether or not s is} \\ & \texttt{the binary code of a~number divisible by 3;} \\ \texttt{else output YES;} \end{array}
```

Such a program P_0 obviously violates the condition supposed for GEN-TEST-DATA. Hence we have to conclude that there is no desired program GEN-TEST-DATA.